

# ON THE CONSISTENCY OF NETWORK DATA WITH PAIRWISE STABILITY

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This paper characterizes the consistency of social network data with pairwise stability, which is a solution concept that in a pairwise stable network, no agents prefer to deviate by forming or dissolving links. I take preferences as unobserved and nonparametric, and seek to characterize the networks that are consistent with pairwise stability. Specifically, given data on a single network, I provide a necessary and sufficient condition for the existence of some preferences that would induce this observed network as pairwise stable. When such preferences exist, I say that the observed network is rationalizable as pairwise stable. Without any restriction on preferences, any network can be rationalized as pairwise stable. Under one assumption that agents who are observed to be similar in the network have similar preferences, I show that an observed network is rationalizable as pairwise stable if and only if it satisfies the Weak Axiom of Revealed Pairwise Stability (WARPS). This result is generalized to include any arbitrary notion of similarity.

As an empirical application, I investigate the extent to which real-world networks are consistent with WARPS. In particular, using the network data collected by [Banerjee et al. \(2013\)](#), I explore how consistency with WARPS is empirically associated with economic outcomes and social characteristics. The main empirical finding is that targeting of nodes that have central positions in social networks to increase the spread of information is more effective when the underlying networks are also more consistent with WARPS.

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## 1. INTRODUCTION

This paper examines theoretically and empirically, the consistency of social and economic networks with pairwise stability, which is a solution concept that in a pairwise stable network, no pair of agents prefer to deviate from the existing network configuration by unilaterally dissolving links, or bilaterally forming links (Jackson and Wolinsky (1996)).

More concretely, suppose we observe a single network of social relationships (who is friends with whom), or a network of financial relationships (who borrows money from whom, and who lends to whom). Now to say that the aforementioned network is pairwise stable, one has to check that agents do not *prefer* to deviate from this network. This requires knowing the preferences of everyone in the network.

In the spirit of classical revealed preference analysis, I assume that preferences are *unobserved* and *nonparametric*<sup>1</sup>, and therefore, consistency of an observed network with pairwise stability requires that there exists some preferences that would induce this network as pairwise stable. This is the first paper to provide a revealed preference characterization of the set of networks that are *rationalizable* as pairwise stable. A network is said to be rationalizable as pairwise stable if and only if there exists preferences that induce the network as pairwise stable.

Without any restriction on preferences, any network is trivially rationalizable as pairwise stable, since we can just say that every agent prefers the observed network to any other networks. However it turns out, under the assumption that agents' preferences are not arbitrarily different from each other, some networks cannot be rationalized as pairwise stable.

More concretely, I impose only one assumption on preferences, which is that agents that are observed to be similar in the network have similar preferences. Now the notion of similarity or "which agents are similar to whom" is represented by a similarity or equivalence relation. Under this restricted heterogeneity assumption, the Weak Axiom of Revealed Pairwise Stability (WARPS) is obtained which characterizes when an observed network is rationalizable as pairwise stable.

The result is then generalized: for any notion of similarity, and under the

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<sup>1</sup>That is, this paper never imposes parametric assumption on preferences

assumption that agents who are similar have similar preferences, I obtain a corresponding revealed preference axiom that characterizes the consistency of social network data with pairwise stability.

Our axiom resembles GARP (Generalized Axiom of Revealed Preference), which characterizes the consistency of consumption data set with utility maximization (Afriat (1967); Samuelson (1938); Varian (1982, 2006)). Just as GARPS is tantamount to an algorithm checking for cycles in revealed preference relations, WARPS here is equivalent to the acyclicity of an appropriately modified profile of revealed preference relations.

For a given notion of similarity, WARPS can be understood as an algorithm that rules out the following revealed preference cycle: take a pair of agents who reveal preferred to form some links (not necessarily with each other), but there exists another pair of similar agents who reveal preferred not to form a link with each other.

One important departure from classical revealed preference theory arises because in networks and other two-sided settings, when a link is not observed between two agents, we are unsure as to which of the two sides blocked and chose not to form a link. Therefore the data do not reveal preferences directly, but induce a multiplicity of possible revealed preference relations.<sup>2</sup>

## 1.1. Application

The second part of the paper is an empirical exercise, the goal is to investigate the extent to which real-world networks are consistent with pairwise stability through consistency with WARPS. In particular, I consider a special implementation of WARPS where similarity is defined using graph isomorphism. That is, similarity relation is defined such that agents are similar when their respective network neighborhoods are isomorphic. This is a conservative notion of similarity that says agents have similar preferences only when their respective network structures are similar. The use of graph iso-

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<sup>2</sup>More concretely, preferences can be thought of as being revealed from network data as follows. If a link is observed between two agents, say that both agents reveal preferred to form this link, conditional on the rest of the network. However when a link is *not* observed between two agents, we can only say that **either** one **or** both of them reveal preferred not to form this link, conditional on the rest of the network.

morphism allows preferences to depend freely on the network structures without specifying how.

WARPS as developed in the theoretical section is either passed or violated by an observed network. For the application, I develop a continuous measure of how consistent an observed network is to WARPS. Moreover, this numerical consistency score is benchmarked with the consistency of a suitably calibrated random (Erdős-Rényi) network with WARPS.<sup>3</sup>

Using these measures, and I demonstrate that the consistency of real-world networks with WARPS matters for economic outcomes. To show this, I use the data set of [Banerjee et al. \(2013\)](#), which consists of a cross-section of 75 villages in rural India. For each village, we observe (i) the social network of the village, (ii) the nodes who were first informed of a new microfinance loan program, as well as (iii) the eventual adoption rate of microfinance.

[Banerjee et al. \(2013\)](#) show that the network centrality of the first-informed agents (injection points) is significantly and positively related to the adoption rate of microfinance. I find that this relationship is significantly more positive when the underlying network is more consistent with WARPS. In showing this, I maintain their exact specification and controls, incorporating only my measures of consistency. Therefore, the consistency of WARPS is new network statistics that has potentially useful application: when a network is more consistent with WARPS, targeting of agents that have central positions in the network to increase the diffusion of information is more effective.

Besides using different measures of consistency as a robustness check, a placebo test is also performed in which a host of other network measures are used in place of our measure of consistency. Like WARPS, these network measures are calculated using only information about the network structure, but unlike WARPS, they were not significant in explaining diffusion. Finally, the result is robust to using a specific network type, the advice network, as the social network of each village (instead of an aggregate of different network types).

To get a sense of magnitudes, in their findings, one standard deviation increase in the diffusion centrality (a measure how central individuals are

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<sup>3</sup>This measure is in line with the notion of predictive success used in the empirical revealed preference literature ([Beatty and Crawford \(2011\)](#); [Selten \(1991\)](#)).

in their social network with regard to spreading information) of injection points would lead to an increase in the eventual adoption rate of 3.9 percentage points. I find that this result depends significantly at how consistent to WARPS the underlying network is. For instance, at the 90th, 75th, and 25th percentile consistency, the corresponding magnitudes are 7.2-8.1, 5.9-6.1, and 1.3-2.0 percentage points increase in the adoption rate.

The intuition for the result is as follows. Networks that are less consistent with WARPS are more difficult to rationalize as pairwise stable. There are only two reasons, either we do not expect pairwise stability to hold, or that there is unobserved heterogeneity not captured by the network structure. For either of these explanations, centrality of injection points *with respect to* the network structure does not correspond to the actual effectiveness of these agents in diffusion. Therefore, the relationship between network centrality and adoption rate is weaker as the underlying network violates WARPS more severely. This is precisely what we see in the data.

## 1.2. Related Literature

This is the first paper to characterize the empirical restriction of pairwise stability in networks. It belongs to the broader agenda of revealed preference in economics, championed by [Afriat \(1967\)](#); [Richter \(1966\)](#); [Samuelson \(1947, 1938\)](#); [Varian \(1982\)](#). They show that when preferences are *unobserved*, the empirical restriction of utility maximization is captured by GARP. Their works led to extensive empirical applications using GARP as a basis for measuring rationality (i.e. consistency with utility maximization). Notably, [Andreoni and Miller \(2002\)](#); [Blundell et al. \(2003\)](#); [Choi et al. \(2014\)](#); [Echenique et al. \(2011\)](#); [Harbaugh et al. \(2001\)](#). For surveys on the subject, see [Crawford and De Rock \(2014\)](#); [Varian \(2006\)](#).

Beyond the setting of individual decision-making, the revealed preference approach has also been applied to other multi-agents (strategic or otherwise) settings. This paper extends to the network setting, although the desideratum is microfounding measures based on revealed preference characterizations, which is not the objective in these papers. For example, [Brown and Matzkin \(1996\)](#) give a revealed preference characterization of Walrasian equilibria; [Carvajal et al. \(2013\)](#) for Cournot equilibria; [Echenique \(2008\)](#); [Echenique et al. \(2013\)](#) for stability of two-sided matching models; [Chambers and Echenique \(2014\)](#) for bargaining equilibria; [Cherchye et al. \(2007,](#)

2011); Chiappori (1988); Sprumont (2000) for collective rationality.

## 2. ILLUSTRATIVE EXAMPLES

Before I formally introduce the general setup, the examples in this section help illustrate the heart of the matter. Consider the social network  $g$  depicted in Figure 1 below. From observing a link (friendship) between agents  $a$  and  $b$ , we say that agents  $a$  and  $b$  both *reveal preferred* the network in Figure 1 to the network in Figure 2. Notationally, the statement “agent  $a$  reveals preferred the network  $g$  to  $g - ab$ ” is represented by  $g \succ_a g - ab$ . Now we can repeat the same step and define revealed preference this way for all pairs of agents who are linked in Figure 1.

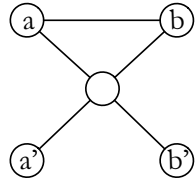


Figure 1: Network  $g$

$\succ_a$   
and  
 $\succ_b$

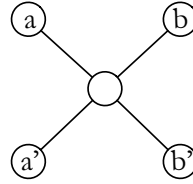


Figure 2: Network  $g - ab$

Now from observing no link between agents  $a'$  and  $b'$ , we say that **either** agent  $a'$  **or**  $b'$  (or both) reveal preferred the network in Figure 3 to the network in Figure 4. While forming a link requires the consent of both sides (bilateral), dissolving a link is unilateral. Therefore when we do not observe a link, we are unsure as to which side prefers not to form the link.

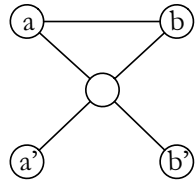


Figure 3: Network  $g$

$\succ_{a'}$   
or  
 $\succ_{b'}$

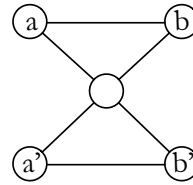


Figure 4: Network  $g + a'b'$

*The question we are interested in is: does there exist some underlying preferences (complete and transitive binary relations) that could induce these revealed prefer-*

ences? If we can find such preferences, then we say the network  $g$  in Figure 1 is rationalizable as pairwise stable.

## 2.1. Network $g$ is not rationalizable as pairwise stable

The network  $g$  in Figure 1 (in fact any network) is rationalizable as pairwise stable.<sup>4</sup> But under the following structure on preferences, it is not rationalizable as pairwise stable. Let an agent's degree in  $g$  be her number of links (friends) in the network  $g$ .

For this example (and this example only), assume that preferences are as follows: there is a ranking of degrees  $d_0, d_1, d_2, \dots$ , such that everyone agrees having  $d_0$  friends is better than  $d_1$  friends, and better than  $d_2$  friends and so on.<sup>5</sup>

From Figures 1 and 2, the degrees of agent  $a$  in the networks  $g$  and  $g - ab$  are 2 and 1 respectively. Therefore the revealed preference  $g \succ_a g - ab$  implies that agent  $a$  reveals preferred to have 2 friends to 1. On the other hand, we see in Figures 3 and 4 that the degrees of agent  $a'$  in the network  $g$  and  $g + a'b'$  are 1 and 2 respectively. However we cannot be sure if agent  $a'$  reveals preferred to have 1 friend than 2 (recall that link deletion is unilateral, and it could have been agent  $b'$  who blocks the formation of the link with  $a'$ ).

Looking again at Figure 3 and 4, the degrees of agent  $b'$  in the network  $g$  and  $g + a'b'$  are also 1 and 2 respectively, just like agent  $a'$ . Therefore even when revealed preferences of agents  $a'$  and  $b'$  are indeterminate, we can still say that there is some agent (either of  $a'$  or  $b'$ ) who reveals preferred 1 friend to 2; yet from the previous paragraph, we found an agent  $a$  who reveals preferred 2 friends to 1. Since agents agree on a preference ranking over degrees, we obtained a **cycle of length 2 in the revealed preference relation** – which implies that network  $g$  in Figure 1 is not rationalizable as pairwise stable.

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<sup>4</sup>To see this, the revealed preference relation we constructed is acyclic, there is therefore an underlying preferences (complete and transitive) that extend or induce them.

<sup>5</sup>For instance, this assumption is satisfied when the utility that agent  $i$  derives from network  $g$  is  $u_i(g) = d_i(g) - cd_i^2(g)$ , where  $d_i(g)$  is the degree of  $i$  in  $g$ , and  $c$  is the cost parameter that is homogeneous among all agents.

## 2.2. Rationalizable networks

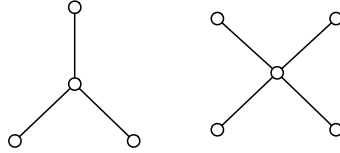


Figure 5: All star networks are rationalizable as pairwise stable.

Given the restrictive assumption on preferences in this example, one can still find networks that are rationalizable as pairwise stable. For example, all star networks with more than 3 agents (Figure 5) are rationalizable as pairwise stable. A preference that rationalizes star networks is as follows: agents prefer fewer friends than more friends below a threshold degree; and agents prefer more friends above a threshold degree. Moreover agents also prefer to have at least one friend to none.

## 2.3. Sufficient condition

Thus far, I have illustrated the intuition behind the necessary condition for a network to be rationalizable as pairwise stable:– check for cycle of length 2 in the revealed preference relations. Is this also sufficient? In particular, one can imagine the possibility of longer cycles, for example, when some agent (reveals) preferred  $d_0$  degrees to  $d_1$ , another prefers  $d_1$  to  $d_2$ , and another prefers  $d_2$  to  $d_0$ .

However a simple argument shows that no cycles of length greater than 2 can exist without nesting a cycle of length 2. Therefore we can *ignore cycle of length greater than 2*. To see this, suppose there is a cycle of length 3, and without loss of generality, say that agent  $i$  (reveals) preferred  $d_0$  number of friends to  $d_0 + 1$  friends. If another agent  $j$  prefers  $d_0 + 1$  to  $d_0$  number of friends, then we obtain a cycle of length 2 in the revealed preference relation. If instead agent  $j$  prefers  $d_0 + 1$  to  $d_0 + 2$  friends, then there can be no agent who reveals preferred  $d_0 + 2$  to  $d_0$ . This is an artifact of pairwise stability, which is a conservative solution concept that requires stable network to be robust to only one-link deviation. As a result, preferences over two networks that differ in more than one link are never revealed.



## 2.4. Modus Operandi

I will demonstrate the main result by applying Theorem 1 to this example, where preference heterogeneity is restricted to degrees (agents that have the same degrees have the same preference). The theorem is applicable generally to other (less restrictive) assumption on preference heterogeneity.

Let  $d_i(g)$  be the degree (or the number of friends) of agent  $i$  in network  $g$ . *The network  $g$  is rationalizable as pairwise stable if and only if for all pairs of agents  $i, j$ , there does not exist a pair of agents  $i', j'$  that are not linked in network  $g$ , such that  $d_i(g) = d_{i'}(g + i'j')$  and  $d_j(g) = d_{j'}(g + i'j')$ .* This condition is essentially an algorithm that decides if a given network data is rationalizable as pairwise stable under the particular restriction on preference heterogeneity of this example. To explain this result further:

Suppose for network  $g$ , the condition above is violated for the pair of agents  $i', j'$ , and the pair of agents  $i, j$ . Denote  $d_i$  for  $d_i(g) = d_{i'}(g + i'j') \geq 1$  and  $d_j$  for  $d_j(g) = d_{j'}(g + i'j') \geq 1$ . Since we observe no link between agents  $i'$  and  $j'$ , **either** agent  $i'$  reveals preferred  $d_i - 1$  to  $d_i$  number of friends, **or** agent  $j'$  reveals preferred  $d_j - 1$  to  $d_j$  number of friends. Otherwise they could have formed a link with each other to achieve  $d_i$  and  $d_j$  number of friends respectively. But we also know that node  $i$  reveals preferred  $d_i$  to  $d_i - 1$  number of friends, otherwise he can unilaterally sever one of his links. Likewise node  $j$  reveals preferred  $d_j$  to  $d_j - 1$  number of friends.

Therefore regardless of the indeterminacy in whether it was agent  $i'$  or  $j'$  who blocked the formation of the link between them, at least one of the following length-2 cycle must be present: (i) agents reveal preferred  $d_i$  to  $d_i - 1$  number of friends, **yet**  $d_i - 1$  to  $d_i$  number of friends, or (ii) they prefer  $d_j$  to  $d_j - 1$  number of friends, **yet**  $d_j - 1$  to  $d_j$  number of friends.

## 3. SETUP

The primitives consist of (i) data, (ii) preferences, and (iii) the notion of pairwise stability.

### 3.1. Data

A network  $g$  is an ordered tuple  $g = (N, E, X)$  where  $N = \{1, 2, \dots, n\}$  is the **set of agents** who can form links with each other. The set  $E$  is the **set of links**. It is a set listing who interacts with whom, also known as an adjacency list. For instance,  $\{i, j\} \in E$  indicates there is a link between agents  $i$  and  $j$ . Implicit is the assumption that the network  $g$  is **undirected**, i.e., if  $i$  is linked to  $j$ , then  $j$  is also linked to  $i$ .

$X = (x_1, \dots, x_n)$  is the vector of agents' characteristics, where  $x_i \in \mathcal{X}_i$  is the characteristics of agent  $i$ . These characteristics are exogenously given and fixed across networks. For instance,  $x_i$  does not depend on who agent  $i$  is linked to, or the network that agent  $i$  is in. As a simplified example of a network, let  $N = \{1, 2, 3\}$ ,  $E = \{\{1, 2\}, \{1, 3\}\}$ , and  $X = (Black, White, Black)$ . The network  $g = (N, E, X)$  then describes a star (or line) network depicted in Figure 6.

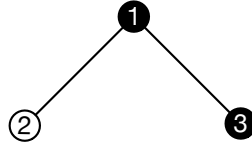


Figure 6:  $g = (N, E, X)$

Let  $\mathcal{G}(N, X)$  be the **set of all networks** with  $N$  as the set of nodes, and  $X$  as the profile of types. Unless stated otherwise, the sets  $N$  and  $X$  are fixed, and we will suppress the dependence of  $\mathcal{G}(N, X)$  on  $N$  and  $X$ .

For ease of exposition, the notation  $ij \in g$  refers to the undirected link  $\{i, j\} \in E$  in the network  $g = (N, E, X)$ . A **non-link**  $ij \notin g$  refers to a pair of agents  $\{i, j\}$  that do not have a link between them in the network  $g = (N, E, X)$ , that is,  $\{i, j\} \notin E$ . In addition, I will use  $g + ij$  to denote the network obtained by adding an undirected link between  $i$  and  $j$  in  $g$ , i.e.,  $g + ij$  is the network  $(N, E \cup \{i, j\}, X)$ . Similarly  $g - ij$  is the network obtained by removing  $ij \in g$ .

### 3.2. Preferences

Agents have preferences over networks. For example, an agent  $i$  might prefer the network  $g$  to  $g - ij$ , denoted as  $g \succeq_i g - ij$ . When  $g \succeq_i g - ij$ , then we also say he prefers to form a link with agent  $j$  in the network  $g$ . In similar veins,  $g \succeq_i g + ij$  means that agent  $i$  prefers not to form a link with agent  $j$  in the network  $g$ .

Recall that  $\mathcal{G}$  is the set of all networks. Now I define agent  $i$ 's **preference space**  $\mathcal{G}_i$  as follows:  $\mathcal{G}_i$  consists of pairs of networks  $(g, g')$  such that  $g$  and  $g'$  are one-link adjacent for agent  $i$ . That is,  $\mathcal{G}_i = \{(g, g') \in \mathcal{G}^2 : \exists j \in N, g = g' + ij \text{ or } g = g' - ij\}$ . That is,  $(g, g') \in \mathcal{G}_i$  if and only if  $g, g'$  differs in one link that involves agent  $i$ .

Agent  $i$ 's preference relation, denoted as  $\succeq_i$ , is formally a subset of  $\mathcal{G}_i$ . It is a binary relation  $\succeq_i \subset \mathcal{G}_i$  that is complete and transitive. For instance, the statement  $(g, g') \in \succeq_i$  is taken to mean  $g \succeq_i g'$ , which says that agent  $i$  weakly prefers the network  $g$  to the network  $g'$ . Moreover by definition of the preference space  $\mathcal{G}_i$ , we can always write  $g \succeq_i g'$  as  $g \succeq_i g \pm ij$ , for some  $j$ . Finally let  $\succeq = (\succeq_1, \dots, \succeq_n)$  be a **profile** of preferences.

### 3.3. Pairwise Stability

In Definition 3.1 below, I introduce pairwise stability, which is the equilibrium condition that no agent wants to dissolve an existing relationship, and no pair of agents want to form a link between them. Pairwise stability was first introduced in Jackson and Wolinsky (1996), and is a weak requirement that individuals have a tendency to form relationships that are mutually beneficial and to drop relationships that are not. Other more restrictive equilibrium concepts exist, for example pairwise Nash equilibrium.<sup>6</sup> Pairwise stability is often considered as a necessary condition for any equilibrium concept in strategic network formation (Calvó-Armengol and İlkılıç (2009)).

**Definition 3.1.** The network  $g \in \mathcal{G}$  is *pairwise stable* with respect to a profile of preferences  $\succeq$  if

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<sup>6</sup>Intuitively, pairwise stable networks are robust to one-link deviations, while pairwise-Nash networks are robust to many-links deletion and single-link creation. See Goyal and Joshi (2003), Goyal (2009) for more.

- 1.) For all  $ij \in g$ ,  $g \succ_i g - ij$  AND  $g \succ_j g - ij$ .
- 2.) For all  $ij \notin g$ ,  $g \succ_i g + ij$  OR  $g \succ_j g + ij$

Implicit in this definition is the idea that while forming links is a collective and bilateral effort of two parties, it only takes one side to dissolve a link.

REMARK 1: Definition 3.1 is the strict version of pairwise stability as defined in Jackson and Wolinsky (1996). In their definition, a link or a non-link can be explained by agents being indifferent to forming a link. This stability concept is too weak: by assuming that agents are always indifferent, any observed networks can be rationalizable as pairwise stable.

### 3.4. Statement of Problem

We can now formally define the statement of problem: what is the necessary and sufficient condition on network  $g$  such that  $g$  is rationalizable as pairwise stable? The notion of rationalizability is more carefully stated below.

**Definition 3.2.** The network data  $g \in \mathcal{G}$  is **rationalizable as pairwise stable** if and only if there exists a profile of preferences  $\succeq = (\succeq_1, \dots, \succeq_n)$  such that  $g$  is pairwise stable with respect to  $\succeq$ .

Without additional assumption on preferences, any network is rationalizable as pairwise stable according to Definition 3.2 – we can just say that agents strictly prefer the observed network to any other networks. The following sections will introduce an assumption that restricts preference heterogeneity: agents that are observed to be similar in the network data have similar preferences. It turns out that under this one assumption, the problem stated here has a rich answer.

### 3.5. Similarity

In order to properly define the restricted heterogeneity assumption, we need a notion of similarity on networks. Only then, we can say that agents who are observed to be similar have similar preferences. Similarity is formalized as an equivalence relation on  $N$  (the set of agents) and  $\mathcal{G}$  (the set of all networks on  $N$ ). The idea that agent  $i$  in network  $g$  is similar to agent  $i'$  in network  $g'$  can be denoted by  $(i, g) \sim (i', g')$ .

**Definition 3.3** (Similarity). A similarity relation  $\sim$ , is an equivalence relation on  $N \times \mathcal{G}$ .

For instance, we can define the similarity relation  $\sim$  to be such that  $(i, g) \sim (i', g')$  if and only if  $\theta_i(g) = \theta_{i'}(g')$ , where  $\theta_i(g)$  is a vector of observed characteristics of agent  $i$  in network  $g$ . The vector  $\theta$  can consist of both exogenous characteristics of agent  $i$  (such as his gender, age, wealth), as well as characteristics that depend on the particular network that agent  $i$  is in. For instance,  $\theta_i(g)$  could tell us the clustering coefficient and network centrality of agent  $i$  in network  $g$ .

### 3.6. Restriction on preference heterogeneity

The final primitive I introduce is the notion of heterogeneity restriction on preferences. Restricted heterogeneity says that agents that are observed to be similar have similar preferences. This assumption on preferences is the only restriction that we will impose on preferences.

Without this assumption, agents' preferences are allowed to be fully heterogeneous, and crucially, heterogeneous in ways that are unobserved to the analyst, so then there always exists some preferences that would rationalize any network data. For the question of what networks can be rationalized as pairwise stable to be meaningful, we need to impose some notion of bounded preference heterogeneity.

Restriction on preference heterogeneity is defined with respect to a given similarity relation, where a similarity relation (Definition 3.3) captures which agents are similar to whom from the analyst's perspective. Therefore, an analyst can propose a notion of similarity, and then heterogeneity of preferences is defined with respect to that.

Definition 3.5 below states the notion of preference heterogeneity restriction. Intuitively, it says that agents that are observed to be similar are deemed to have the same preferences. We will always specify similarity in terms of observable characteristics of the agents, therefore this assumption is tantamount to assuming that there is *no unobserved heterogeneity in preferences*. In particular, names or identities of the agents do not matter.

The restricted heterogeneity assumption we impose on preferences is weak in the sense that it does not restrict preferences for networks that are off-

path. Off-path networks are those that are more than one-link different from the observed network. Therefore we are completely agnostic about preferences of agents when they are off-path from the observed network.

**Definition 3.4.** Let  $g$  be an observed network. The set of **on-path networks** with respect to  $g$ ,  $\mathcal{G}(g)$ , consists of networks that are one-link adjacent to  $g$ . That is,  $\mathcal{G}(g) = \{g' \in \mathcal{G} : \exists ij \in g \text{ or } \exists i'j' \notin g \text{ such that } g' = g - ij \text{ or } g' = g + i'j'\}$ . In words, the set of on-path networks,  $\mathcal{G}(g)$ , consists of networks that are obtainable from  $g$  through a single deletion or addition of link. Therefore when observing the network  $g$ , pairwise stability only reveals preferences for on-path networks with respect to  $g$ , i.e. those networks in  $\mathcal{G}(g)$ .

**Definition 3.5.** Let  $g^0$  be the observed network. The preference profile  $\succeq$  is **heterogeneous with respect to the similarity relation**  $\sim$  at  $g^0$  if and only for all on-path networks:  $g, g - ij$  and  $g', g' - i'j'$ :  
 *$(i, g) \sim (i', g')$  and  $(i, g - ij) \sim (i', g' - i'j')$  implies that agent  $i$  must rank the networks  $g$  and  $g - ij$  the same way as agent  $i'$  ranks the networks  $g'$  and  $g' - i'j'$ .*<sup>7</sup>

Imagine an agent  $i$  choosing between two alternatives,  $(i, g)$  and  $(i, g - ij)$ , that is, he chooses whether to form a link with agent  $j$  in network  $g$ . Now agent  $i'$  is also choosing between two alternatives,  $(i', g')$  and  $(i', g' - i'j')$ . Heterogeneity with respect to  $\sim$  says that because agents  $i, i'$  are facing choice problems that are observationally similar according to  $\sim$ , they must have the same preference (ordinal ranking) in these two parallel choice problems. Whenever agent  $i$  prefers to form a link with  $j$  in  $g$ , then agent  $i'$  prefers to form a link with  $j'$  in  $g'$ .

In fact, the extensive example discussed in the last section imposes the restricted heterogeneity assumption where  $\sim$  is defined to be  $(i, g) \sim (i', g')$  whenever  $d_i(g) = d_{i'}(g')$ , the number of links (degree) of  $i$  in  $g$  is denoted by  $d_i(g)$ . Note that in this example,  $(i, g) \sim (i', g') \iff (i, g - ij) \sim (i', g' - i'j')$

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<sup>7</sup>More precisely, either (i)  $g \succ_i g - ij$  and  $g' \succ_{i'} g' - i'j'$ , or (ii)  $g \prec_i g - ij$  and  $g' \prec_{i'} g' - i'j'$ .

#### 4. MAIN RESULT

This section presents testable conditions that characterize what networks are rationalizable as pairwise stable. Only one restriction is imposed on preferences, that is preferences are heterogeneous with respect to a given similarity relation  $\sim$ , i.e. agents that are similar according to the similarity relation  $\sim$  have similar preferences.

The following axiom, the Weak Axiom of Revealed Pairwise Stability, characterizes the testable implication of pairwise stability (under the assumption of restricted heterogeneity).

**Definition 4.1** (WARPS). Let  $\sim$  be an equivalence relation on  $N \times \mathcal{G}$ . The network  $g$  satisfies the Weak Axiom of Revealed Pairwise Stability (WARPS) for  $\sim$  if and only if for all pairs of links  $(i, x), (j, y) \in g$ , there does not exist a non-link  $(i', j') \notin g$  such that:

- (i)  $(i, g) \sim (i', g + i'j')$  and  $(i, g - ix) \sim (i', g)$
- (ii)  $(j, g) \sim (j', g + i'j')$  and  $(j, g - jy) \sim (j', g)$

**THEOREM 1.** Let  $g$  be an observed network, and  $\sim$  be an equivalence relation on  $N \times \mathcal{G}$ . Assume that preferences are heterogeneous with respect to the similarity relation  $\sim$  at  $g$ . An observed network  $g$  is rationalizable as pairwise stable if and only if  $g$  satisfies WARPS for  $\sim$  (Definition 4.1).

Violation of WARPS by the network  $g$  implies that there is a link that should be formed (say between agents  $i'$  and  $j'$ ), but it is not observed in the data  $g$ . As a consequence, there does not exist preferences that would rationalize  $g$  as pairwise stable. How do we know that the a link should be formed between agents  $i'$  and  $j'$ ? From observing the revealed preference of other similar agents (say  $i$  and  $j$ ). The choices of these similar agents reveal that they would like to form links  $((i, x)$  and  $(j, y)$  respectively) in  $g$ , which implies that agents  $i', j'$  should also form the non-link  $(i', j') \notin g$ . Preference heterogeneity with respect to  $\sim$  allows us to identify that agents  $i', j'$  have similar preferences to  $i, j'$  respectively.

A formal argument based on the above paragraph is used to prove that WARPS is necessary (Section 4.1 below). From this result, we can see that when too few agents are similar according to  $\sim$ , then WARPS is always satisfied by all networks. In particular, when  $\sim$  is defined such that  $(i, g) \sim$

$(i', g')$  only if the networks  $g$  and  $g'$  have the same number of links, then WARPS is always satisfied, and all networks are rationalizable as pairwise stable. When WARPS has no power, we would need to impose a weaker heterogeneity restriction via  $\sim$ .

WARPS here is tantamount to an algorithm checking for cycles in an appropriately constructed profile of revealed preference relations. Existence of transitive and complete preference relations can then be found by extending these acyclic revealed preference relations. This is the technique used to prove the sufficiency of WARPS, which is the more difficult direction. The adjective 'weak' in WARPS is used because WARPS checks for cycles of length 2, which is reminiscent of the Weak Axiom of Revealed Preference (WARP) in consumer theory.

#### 4.1. Outline of Proof

Instead of working with the similarity relation  $\sim$  as in Definition 3.3, it will be more convenient to work with the following induced equivalence relation  $\sim^*$ .

**Definition 4.2.** Let  $\sim^*$  be an equivalence relation on  $N \times \mathcal{G}^2$  induced by the similarity relation  $\sim$  such that  $(i, g, g - ij) \sim^* (i', g', g' - i'j')$  if and only if  $(i, g) \sim (i', g')$  and  $(i, g - ij) \sim (i', g' - i'j')$

Following this new definition of similarity relation, a preference profile  $\succeq$  is heterogeneous with respect to the similarity relation  $\sim^*$  if and only if  $(i, g, g - ij) \sim^* (i', g', g' - i'j')$  implies that agent  $i$  must rank the networks  $g$  and  $g - ij$  the same way as agent  $i'$  ranks the networks  $g'$  and  $g' - i'j'$ .

##### 4.1.1. Necessity

We will now argue that WARPS is necessary. So suppose that WARPS is violated for a network  $g$ , but there exists some preferences that rationalize  $g$  as pairwise stable. The next two paragraphs will demonstrate a contradiction. Begin with the fact that there is no link between agents  $i'$  and  $j'$ . From pairwise stability, we can only infer that *either or both* of  $i'$  and  $j'$  prefers  $g$  to  $g + i'j'$ .

Consider the first possibility that it is agent  $i'$  who prefers not to form the



link  $i'j'$ . Now from observing the link  $(i, x)$  in the data, it must be that  $g \succ_i g - ix$ . Moreover  $(i, x)$  is such that  $(i, g, g - ix) \sim^* (i', g + i'j', g)$ . We can then infer that agents  $i, i'$  have similar preference, and in particular,  $g \succ_i g - ix$  implies that  $g + i'j' \succ_{i'} g$ . This contradicts the assumption that agent  $i'$  prefers  $g$  to  $g + i'j'$ .

Consider the second possibility that it is agent  $j'$  who prefers not to form the link  $i'j'$ . Now from observing the link  $(j, y)$  in the data, it must be that  $g \succ_j g - jy$ . Moreover  $(j, y)$  is such that  $(j, g, g - jy) \sim (j', g + i'j', g)$ . We can then deduce that agents  $j, j'$  have similar preference, and in particular,  $g \succ_j g - jy$  implies that  $g + i'j' \succ_{j'} g$ . Again this contradicts the assumption that agent  $j'$  prefers  $g$  to  $g + i'j'$ .

#### 4.1.2. Sufficiency

To show that WARPS is sufficient, we will utilize Lemma 1, which says if we can elicit a profile of revealed preference relations from the observed network  $g$  such that it is acyclic when extended according to the similarity relation  $\sim$ , then there exist complete and transitive preferences that rationalize  $g$  as pairwise stable (and satisfy restricted heterogeneity). The main machinery is Algorithms 1 and 2, which constructively demonstrate a profile of revealed preference relation whose extension according to  $\sim$  is acyclic.

First, we will define how to elicit revealed preference relations from the data. If a link between agents  $i$  and  $j$  is observed in the network  $g$ , then we say that agents  $i$  and  $j$  both *reveal preferred* to have this link to not having this link in  $g$ . On the other hand, if we do not observe a link between agents  $i$  and  $j$  in the network  $g$ , then it is revealed that *either* agent  $i$  *or* agent  $j$  prefers not to form this link in  $g$ . More formally,

**Definition 4.3.** For a given network  $g$ , let  $\succ^R = \prod_{i \in N} \succ_i^R$  be a profile of **revealed preference relations** defined as follows:

- 1.) For all  $ij \in g$ ,  $g \succ_i^R g - ij$  AND  $g \succ_j^R g - ij$ .
- 2.) For all  $ij \notin g$ ,  $g \succ_i^R g + ij$  OR  $g \succ_j^R g + ij$

This definition does not pin down a unique revealed preference relation (due to part 2.) of the definition). Now let  $\mathcal{R}(g)$  be the set of all such profiles of revealed preference relations derived via Definition 4.3 above. The set

$\mathcal{R}(g)$  contains all the possible revealed preference relations that are induced by the data  $g$ .

**Definition 4.4.** The network  $g \in \mathcal{G}$  is **rationalizable** as pairwise stable if and only if there exists a profile of strict preferences that extends some profile of revealed preference relations in  $\mathcal{R}(g)$ .

In classical revealed preference theory, the revealed preference relation is uniquely defined. In our setup, rationalization only requires that we can find just one revealed preference relation in  $\mathcal{R}(g)$  that can be extended into a complete and transitive relations.

**Definition 4.5.** Let  $\succ = \prod_{i \in N} \succ_i$  be a profile of binary relations. The extension of  $\succ^R$  according to the similarity relation  $\sim$  is a profile of binary relations  $R = (R_1, \dots, R_n)$  such that if  $g \succ_i g - ij$ ,  $(i, g) \sim (i', g')$  and  $(i, g - ij) \sim (i', g' - i'j')$ , then  $g' R_i g' - i'j'$

**Definition 4.6.** The profile of strict (asymmetric) binary relation  $\succ^R = (\succ_1^R, \dots, \succ_n^R)$  is acyclic if for all  $i \in N$ , there is no sequence  $x_1, x_2, \dots, x_L$  such that  $x_1 \succ_i^R x_2 \succ_i^R \dots x_L$  and  $x_L \succ_i^R x_1$ .

**Lemma 1.** *Assume that preferences are heterogeneous with respect to the similarity relation  $\sim$ . The network  $g$  is rationalizable as pairwise stable if and only if there exists a profile of revealed preference relations  $\succ^R \in \mathcal{R}(g)$  such that the extension of  $\succ^R$  according to  $\sim$  is acyclic.*

Lemma 1 is an application of a version of *Szpilrajn extension theorem*. If a revealed preference relation is acyclic, then we can extend it to a complete and transitive preference relation. Since we also want the complete and transitive preferences to satisfy heterogeneity with respect to  $\sim$ , the revealed preference relations are first extended according to  $\sim$ , and it is the acyclicity of this  $\sim$ -extended relations that matter.

The profile of revealed preference relations itself is always acyclic by definition, but its  $\sim$ -extension might not be. In fact, we now show that when WARPS for  $\sim$  is satisfied, there exists a profile of revealed preference relation whose  $\sim$ -extension is acyclic, where  $\sim$  is some similarity relation. Throughout the rest of this section, it will be maintained that (a)  $g$  denotes the observed network, (b)  $g$  satisfies WARPS for  $\sim$ , and (c)  $\sim$  is some simi-

larity relation.

The application of Algorithms 1 and 2 successively will demonstrate existence constructively. The input to the first algorithm is a special profile of relations that are elicited from  $g$  by assuming that link deletion is bilateral. This revealed preference relation is obtained by inferring that when a link is not observed between two individuals, *both* sides do not want to form this link. See Definition 4.7 below.

**Definition 4.7.** For an observed network  $g$ , let  $\succ^0 = \prod_{i \in N} \succ_i^0$  be the unique profile of revealed preference relations defined as follows:

- 1.) For all  $ij \in g$ ,  $g \succ_i^0 g - ij$  AND  $g \succ_j^0 g - ij$
- 2.) For all  $ij \notin g$ ,  $g \succ_i^0 g + ij$  AND  $g \succ_j^0 g + ij$

Let  $R$  be an  $\sim$ -extension of some  $\succ \in \mathcal{R}(g)$ . To understand the following algorithms, observe that all possible cycles in  $R$  must take the following form: either  $(g R_i g - ix$  and  $g - ix R_i g)$ , or  $(g R_i g + ix$  and  $g + ix R_i g)$ . This is because in Definition 3.5, restricted heterogeneity is defined only locally on  $g$ , it only restricts preferences for on-path networks, hence restricted hetero-

generity only causes cycles in  $R$  that involve on-path networks.

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**Algorithm 1:** Eliminate cycles I

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**input** :  $\succ$  as the profile of revealed preference relations constructed in Definition 4.7  
**output**:  $\succ^1$ , a profile of binary relations

- 1 **while** there is a cycle in  $R$  of the form  $(g R_i g - ix$  and  $g - ix R_i g)$ , where  $R$  is the extension of  $\succ$  according to  $\sim$ , **do**
- 2     identify a cycle of the form  $(g R_i g - ix$  and  $g - ix R_i g)$ . The relation  $g - ix R_i g$  is true if and only if there is a sequence of agents  $(i_0, i_2, \dots, i_{M-1})$  and a sequence of non-links  $(i_0, j_0), \dots, (i_{M-1}, j_{M-1}) \notin g$  such that for all  $m = 1, \dots, M - 1$ ,  $(i, g, g - ix) \sim^* (i_m, g + i_m j_m, g)$  and  $g \succ_{i_m} g + i_m j_m$ ;
- 3     **for**  $m = 1, \dots, M$  **do**
- 4         **replace**  $(g \succ_{i_m} g + i_m j_m)$  with  $(g \prec_{i_m} g + i_m j_m)$ ;
- 5         **for each**  $(i', j') \notin g$  such that  $(i', g, g + i' j') \sim^* (i_m, g, g + i_m j_m)$  **do**
- 6             **replace**  $(g \succ_{i'} g + i' j')$  with  $(g \prec_{i'} g + i' j')$ ;
- 7         **end**
- 8     **end**
- 9 **end**

---

In Lemma 2 below, we will show that Algorithm 1 terminates in finite steps, and crucially, we will claim that the output  $\succ^1$  from the algorithm belongs to the set  $\mathcal{R}(g)$ , i.e. it is a profile of revealed preference relations.

**Lemma 2.** *Algorithm 1 terminates in finite steps. The output of Algorithm 1 is a profile of revealed preference relations, i.e.  $\succ^1 \in \mathcal{R}(g)$ .*

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**Algorithm 2:** Eliminate cycles II
 

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**input** :  $\succ$  as  $\succ^1$ , the profile of revealed preference relations from Algorithm 1;  
**output**:  $\succ^2$ , a profile of binary relations;

- 1 **while** there is a cycle in  $R$  of the form  $(g R_i g + ix$  and  $g + ix R_i g)$ , where  $R$  be the extension of  $\succ$  according to  $\sim$ , **do**
- 2     identify a cycle of the form  $(g R_i g + ix$  and  $g + ix R_i g)$ ;
- 3     **if**  $(g \succ_i g + ix)$  is true **then**
- 4         **replace**  $(g \succ_i g + ix)$  with  $(g \prec_i g + ix)$ ;
- 5         **for each**  $(i', j') \notin g$  such that  $(i', g, g + i'j') \sim^* (i, g, g + ix)$ , and  $(g \succ_{i'} g + i'j')$  **do**
- 6             **replace**  $(g \succ_{i'} g + i'j')$  with  $(g \prec_{i'} g + i'j')$ ;
- 7         **end**
- 8     **end**
- 9 **end**

---

**Lemma 3.** *Algorithm 2 terminates in finite steps. The output  $\succ^2$  is a profile of revealed preference relations, i.e.  $\succ^2 \in \mathcal{R}(g)$ . Moreover, the extension of  $\succ^2$  according to  $\sim$  is acyclic.*

What we have shown is whenever WARPS for  $\sim$  is satisfied, we can construct (through Algorithms 1 and 2), a profile of revealed preference relations from the observed network, such that the extension according to  $\sim$  is acyclic. By Lemma 1, this extended profile can be further extended into a complete and transitive preference relations that satisfy heterogeneity with respect to  $\sim$ .

## 5. UNOBSERVED HETEROGENEITY

In this section, agents are allowed to have unobserved characteristics. When there is unbounded unobserved heterogeneity, every agent can have different preferences, and as a result, we can rationalize any network as pairwise stable (WARPS is always satisfied). However we can ask if under limited but reasonable amount of unobserved heterogeneity, a network would be rationalizable as pairwise stable.

The main result of this section is the formulation of WARPS with randomness (Proposition 1), which characterizes the likelihood that a network is rationalizable as pairwise stable, given some heterogeneity restriction on *both* observed and unobserved characteristics. The source of randomness is driven by the distribution of unobserved characteristics. The second result then derives the expected minimum number of types of agents that are needed to rationalize the network data as pairwise stable. Here, types are endogenous partition of the agents' characteristics (both observed and unobserved).<sup>8</sup>

The analysis with unobserved characteristics is based on the previous sections, the main difference is that the similarity relation,  $\sim$ , between agents is now probabilistic, to reflect our uncertainty about the unobserved characteristics agents have. As a result, whether or not an observed network satisfies WARPS is now a random variable, whose realization depends on the draw from the distribution of unobserved characteristics. This random version of WARPS is what characterizes the likelihood that a network is rationalizable as pairwise stable. For instance, we can then ask what is the probability that WARPS is rejected.

Let  $\epsilon_i$  be a vector of (exogenous) unobserved characteristics of agent  $i$ . Agent  $i$ 's true vector of characteristics is  $\tau_i(g)$ , which can depend on the network  $g$ , that agent  $i$  is in. We will assume that  $\tau_i(g) = \theta_i(g) + \epsilon_i$ , where  $\theta_i(g)$  is the vector of observed characteristics of agent  $i$  in network  $g$ . The vector  $\theta_i(g)$  can contain both exogenous characteristics of  $i$ , and the characteristics of  $i$  that is endogenous to the network  $g$ . For instance,  $\theta_i(g)$  could have contain rows indicating the age and wealth of agent  $i$ , and it could contain rows indicating the network centrality and the degree of agent  $i$  in network  $g$ .

**Definition 5.1** (Similarity relation with unobserved characteristics). Let  $\sim$  be an equivalence relation on  $N \times \mathcal{G}$  such that  $(i, g) \sim (i', g')$  if and only if  $|\tau_i(g) - \tau_{i'}(g')| < \delta$ , where  $\delta \in \mathbb{R}_+^L$  and  $\tau_i(g) = \theta_i(g) + \epsilon_i$ . Now  $\theta : N \times \mathcal{G} \rightarrow \mathbb{R}^L$  is an  $L$ -dimensional vector of observed characteristics. For all  $i \in N$ ,  $\epsilon_i$  is an  $L$ -dimensional vector of random variable.  $(\epsilon_1, \dots, \epsilon_n)$  is jointly distributed according to the distribution  $F(\cdot)$ . For simplicity, we will further assume that  $(\epsilon_i)_{i \in N}$  are i.i.d.

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<sup>8</sup>A paper that is similar in spirit is Crawford and Pendakur (2013), which asks how many types of consumers are there that would rationalize a consumption data set.

**PROPOSITION 1** (Random WARPS). *Let  $\sim$  be a random similarity relation defined above. Let  $g$  be the observed network and assume that preferences are heterogeneous with respect to  $\sim$  at  $g$ . The likelihood that a network  $g$  is rationalizable as pairwise stable is given by  $\Pr[T(g, \delta) = 1]$  in Equation (5.1) below, where the expectation is taken with respect to  $F$ , the joint probability distribution of  $\epsilon = (\epsilon_1, \dots, \epsilon_n)$ .*

$$(5.1) \quad \mathbb{E} \left[ \max_{(i,x),(j,y) \in g, (i',j') \notin g} \left\{ \mathbb{1} \left[ |\tau_i(g) - \tau_{i'}(g + i'j')| < \delta \right] \cdot \mathbb{1} \left[ |\tau_i(g - ix) - \tau_{i'}(g)| < \delta \right] \right. \right. \\ \left. \left. \cdots \mathbb{1} \left[ |\tau_j(g) - \tau_{j'}(g + i'j')| < \delta \right] \cdot \mathbb{1} \left[ |\tau_j(g - jy) - \tau_{j'}(g)| < \delta \right] \right\} \right]$$

**PROOF:** The idea is to rewrite WARPS and use Theorem 1. WARPS in Definition (4.1) can be represented by an indicator function  $T(g, \sim) \in \{0, 1\}$ . When  $T(g, \sim) = 0$ , WARPS is satisfied for  $g$ , and when  $T(g, \sim) = 1$ , WARPS is rejected for  $g$ .

$$T(g, \sim) = \max_{(i,x),(j,y) \in g, (i',j') \notin g} \left\{ \mathbb{1} [(i, g) \sim (i', g + i'j')] \cdot \mathbb{1} [(i, g - ix) \sim (i', g)] \right. \\ \left. \cdots \mathbb{1} [(j, g) \sim (j', g + i'j')] \cdot \mathbb{1} [(j, g - jy) \sim (j', g)] \right\}$$

Theorem 1 says that when there is no unobserved characteristics, a network  $g$  is rationalizable as pairwise stable (with preferences that are heterogeneous with respect to  $\sim$ ) if and only if  $T(g, \sim) = 1$ .

When there is unobserved characteristics,  $T(g, \sim)$  becomes a random variable. For each draw of unobserved heterogeneity  $(\epsilon_1, \dots, \epsilon_n)$  from  $F(\cdot)$ , a value of  $T(g, \sim) \in \{0, 1\}$  is realized. In particular, the likelihood that  $g$  is rationalizable as pairwise stable is the probability that  $\Pr[T(g, \sim) = 1]$ , which is just  $\mathbb{E}[T(g, \sim)] = \int T(g, \sim) dF$ .

□

Provided that we specify what the threshold  $\delta$  is, the quantity of interest,  $\Pr[T(g, \delta) = 1]$  in Equation (5.1) can be computed by simulation. For instance, assuming  $(\epsilon_i)_{i \in N}$  are independent, we can set  $\delta = 0.1 \text{diag}(\Sigma)$  where  $\text{Var}(\epsilon_i - \epsilon_{i'}) = \Sigma$ , and  $\text{diag}(\Sigma)$  is the vector of the diagonals of  $\Sigma$ .

We can also ask what is the largest  $\delta$  such that the data  $g$  is rationalizable at least  $1 - \alpha\%$  of the time. When  $\delta$  is small enough, WARPS is never rejected and any network is rationalizable. We can think of  $\delta$  as inversely proportional to how much heterogeneity there is among agents. If we need very large  $\delta$  to achieve  $1 - \mathbb{E}[T(g, \delta)] \geq 1 - \alpha$ , or  $\alpha \geq \mathbb{E}[T(g, \delta)]$ , then it means that we even under limited heterogeneity, the network  $g$  can be rationalized as pairwise stable.

To elaborate on this, let  $\delta^*$  as defined in Equation (5.2) below. Equation (5.2) says that we want to minimize the objective function,  $\Pr[|\tau_i(g) - \tau_{i'}(g)| < \delta]$ , which quantifies how likely are two agents similar to each other under some  $\delta$ . The constraint is that WARPS is rejected less than  $\alpha\%$  of the time. Therefore,  $N^*$  in Equation (5.3) can be interpreted as follows: we need at least  $N^*$  number of types in order to confidently rationalize the data  $g$  as pairwise stable at the  $\alpha\%$  level.

**Definition 5.2.** The expected minimum number of types of agents needed to rationalize the network  $g$  as pairwise stable at least  $1 - \alpha\%$  of the time is  $N^*$ , which solves the following two equations. The function  $\mathbb{E}[T(g, \delta)]$  depends on  $\delta$  and is given in Equation (5.1).

$$(5.2) \quad \delta^* = \operatorname{argmax}_{\delta \in \mathbb{R}_+^L} \left\{ \Pr[|\tau_i(g) - \tau_{i'}(g)| < \delta] \right\} \text{ subject to } \alpha \geq \mathbb{E}[T(g, \delta)]$$

$$(5.3) \quad N^* = \frac{1}{\Pr[|\tau_i(g) - \tau_{i'}(g)| < \delta^*]}$$

The number  $N^*$  informs us the difficulty of rationalizing a network as pairwise stable. It tells us how much preference heterogeneity we need to rationalize a network as pairwise stable. If  $N^*$  is large relative to the number of agents in the network, then rationalizing the network as pairwise stable comes at a substantial cost of relying on the assumption that agents have myriad of different preferences. All else the same, we desire a smaller  $N^*$  relative to the size of the network, because then the data is rationalizable with a simpler model of preferences.



## 6. ADDITIONAL PROOFS

### 6.1. Proof of Lemma 2

PROOF OF LEMMA 2: We want to show that the output of the algorithm is a profile of revealed preference relations. This consists of showing that WARPS implies the algorithm does not result in a profile of relations where there exists a pair of non-link  $i'j' \notin g$  such that  $(g + i'j' \succ_{i'} g)$  and  $(g + i'j' \succ_{j'} g)$ . That is, this would imply that both agents  $i', j'$  prefer to form the non-link  $i'j' \notin g$ , which is not consistent with Definition 4.3. If for some  $(i', j') \notin g$  the algorithm changes  $(g \succ_{i'} g + i'j')$  to  $(g + i'j' \succ_{i'} g)$ , and  $(g \succ_{j'} g + i'j')$  to  $(g + i'j' \succ_{j'} g)$ , then there must exist links  $(i, x), (j, y) \in g$  such that  $(i, g, g - ix) \sim^* (i', g + i'j', g)$  and  $(j, g, g - jy) \sim^* (j', g + i'j', g)$ . This is precisely ruled out by WARPS.

Moreover the algorithm must terminate in finite steps because at each iteration of the while-loop, the relation  $g - ix R_i g$  is reversed, and hence one cycle of the form  $(g R_i g + ix$  and  $g + ix R_i g)$  is eliminated, and no new cycle of that form is introduced. Therefore each iteration of the while-loop strictly decreases the number of cycles in  $R$ , and the algorithm must terminate in finite steps.  $\square$

### 6.2. Proof of Lemma 3

PROOF OF LEMMA 3: We will first prove by induction that whenever there is a cycle of the form  $(g R_i g + ix$  and  $g + ix R_i g)$ , there is only one possibility:  $g \succ_i g + ix$ , but there is a link  $(p, q) \in g$  such that  $(p, g, g - pq) \sim^* (i, g + ix, g)$ . The two other possibilities are (i)  $g \succ_i g + ix$ , but there is a non-link  $(i', j') \notin g$  such that  $g + i'j' \succ_{i'} g$ , and  $(i', g + i'j', g) \sim^* (i, g + ix, g)$ , or (ii)  $g + ix \succ_i g$ , but there is a non-link  $(i', j') \notin g$  such that  $g \succ_{i'} g + i'j'$ , and  $(i', g + i'j', g) \sim^* (i, g + ix, g)$ .

We assume that this hypothesis holds for the 1st to  $t$ -th iterations of the while loop. Assume we are at the  $t + 1$ -th iteration. Cases (i) and (ii) are directly ruled out by the for-loop in Lines (5) and (5) of Algorithms 1 and 2 respectively (similarly for base case of the induction).

At each iteration of the while loop, the cycle  $(g R_i g + ix$  and  $g + ix R_i g)$  is destroyed because  $g \succ_i g + ix$  is switched to  $g \prec_i g + ix$ . Moreover, no new

cycle will be introduced. If there is a new cycle, then it would be because Case (i) or (ii) arises. Therefore the algorithm must terminate.

Finally we prove that the output of the algorithm  $\succ^2$  is a profile of revealed preference relations, i.e. it belongs to  $\mathcal{R}(g)$ . The proof consists of showing WARPS implies that the algorithm does not result in a profile of relations such that there exists a pair of non-link  $i'j' \notin g$  such that  $(g + i'j' \succ_{i'} g)$  and  $(g + i'j' \succ_{j'} g)$ . That is, both agents  $i', j'$  prefer to form the non-link  $i'j' \notin g$ , which is not consistent with Definition 4.3. If the algorithm changes  $(g \succ_{i'} g + i'j')$  to  $(g + i'j' \succ_{i'} g)$ , and  $(g \succ_{j'} g + i'j')$  to  $(g + i'j' \succ_{j'} g)$ , then there exists links  $(i, x), (j, y)$  such that  $(i, g, g - ix) \sim^* (i', g + i'j', g)$  and  $(j, g, g - jy) \sim^* (j', g + i'j', g)$ . This is precisely ruled out by WARPS.  $\square$

## 7. APPLICATION

In this section, I will consider an application of the theoretical result obtained in Section 4. Firstly, a special case of the Weak Axiom of Revealed Preference (WARPS) is derived. This special case is derived by imposing a notion of similarity where agents are similar whenever their respective network neighborhoods are isomorphic. Under the assumption that preferences are restricted to be heterogeneous with respect to this similarity relation (i.e. observationally similar agents have similar preferences), a special case of WARPS is obtained using Theorem 4.1.

Secondly, I propose a continuous measure of consistency with WARPS, and empirically examine the consistency of real-world networks with WARPS. In particular, I use the data set (Banerjee et al. (2013)) consisting of a cross-section of 75 villages in rural India. For each village, we observe (i) the social network of the village, (ii) the nodes who were first informed of a new microfinance loan program, as well as (iii) the eventual adoption rate of microfinance. The main empirical finding is that: targeting of nodes that have central positions in the underlying social network to increase the spread of a new microfinance program is more effective when the network is more consistent with WARPS.

### 7.1. Graph theoretic definitions

**Definition 7.1** (Isomorphism). Consider networks  $g, g' \in \mathcal{G}$ ,  $g$  and  $g'$  are isomorphic if and only if there is a bijection  $\sigma : N \rightarrow N$  such that  $ij \in g \iff \sigma(i)\sigma(j) \in g'$ .

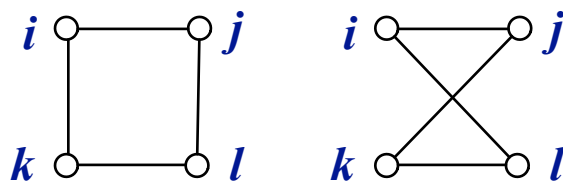


Figure 7: The left network:  $g = \{ij, ik, jl, kl\}$ , and the right network:  $g' = \{ij, il, jk, kl\}$  are isomorphic. Both networks represent the same network structure.

Two networks are isomorphic if they have the same network structure. That is, we can relabel and permute the identities of the nodes in one network to obtain another isomorphic network. For example, consider the network  $g = \{ij, ik, jl, kl\}$  depicted in the left side of Figure 7 below. Consider a permutation  $\sigma$  on the labels of the nodes such that  $\sigma(k) = l$ ,  $\sigma(l) = k$ ,  $\sigma(i) = i$  and  $\sigma(j) = j$ . Then, we can see that  $\sigma(g) = \{ij, il, jk, lk\} = g'$ .

**Definition 7.2** ( $k$ -neighborhoods). The **neighborhood** of node  $i$  in network  $g$  is the set of nodes that  $i$  is linked to. While the  $k$ -neighborhood of node  $i$  in network  $g$  is the set of all nodes that are distance no more than  $k$  from  $i$ . Now the  $k$ -**neighborhood network** of node  $i$  in network  $g$  is the subgraph induced by the nodes in the  $k$ -neighborhood of  $i$ . For simplicity,  $k$ -neighborhood will be understood as the  $k$ -neighborhood *network*.

More formally,  $N_i(g) = \{i\} \cup \{j : ij \in g\}$  is the neighborhood of node  $i$  in  $g$ . The  $k$ -neighborhood is recursively defined as  $N_i^k(g) = N_i(g) \cup \left( \bigcup_{j \in N_i(g)} N_j^{k-1}(g) \right)$ . The  $k$ -neighborhood graph of node  $i$  in  $g$ , denoted as  $g(N_i^k(g))$ , is the *subgraph* of  $g$  with  $N_i^k(g)$  as the set of nodes such that  $i'j' \in g(N_i^k(g))$  if and only  $i', j' \in N_i^k(g)$  and  $i'j' \in g$ .

## 7.2. $k$ -WARPS

A specific form of the Weak Axiom of Revealed Pairwise Stability (WARPS) is provided here. The main result of the previous section, Theorem 1, is applied to obtain Corollary 1 here.

**Definition 7.3** ( $k$ -similar). Let  $\sim_k$  be an equivalence relation on  $N \times \mathcal{G}$  such that  $(i, g) \sim_k (i', g')$  (read as agent  $i$  in network  $g$  is **k-similar** to agent  $i'$  in network  $g'$ ) if and only if there exists an isomorphism  $\sigma$  from the  $k$ -neighborhood network of  $i$  to that of  $i'$  such that  $\sigma(i) = i'$ .

When the similarity relation  $\sim_k$  is as in Definition 7.3, we can simplify the statement of WARPS through the following observation: whenever  $(i, g) \sim_k (i', g + i'j')$ , then there is a link  $ix \in g$  such that  $(i, g - ix) \sim_k (i', g)$ .

**Definition 7.4** ( $k$ -WARPS). Let  $\sim_k$  be a similarity relation from Definition 7.3. The network  $g$  satisfies  $k$ -WARPS if and only if for all pairs of agents  $(i, j)$ , there does not exist a non-link  $i'j' \notin g$  such that:

- (i)  $(i, g) \sim_k (i', g + i'j')$
- (ii)  $(j, g) \sim_k (j', g + i'j')$

**Corollary 1:** Let  $g$  be the observed network, and let  $\sim_k$  be a similarity relation in Definition 7.3. Assume that preferences are heterogeneous with respect to  $\sim_k$  at  $g$ . The network  $g$  is rationalizable as pairwise stable if and only if  $g$  satisfies  $k$ -WARPS.

### 7.3. Illustration

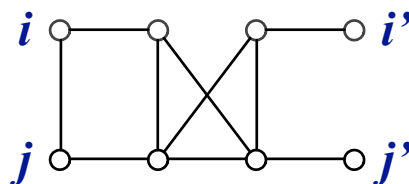
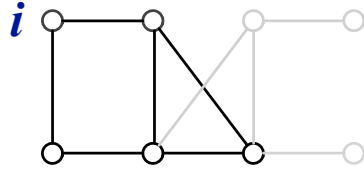
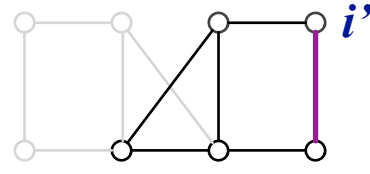


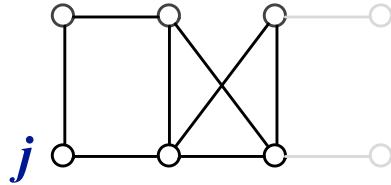
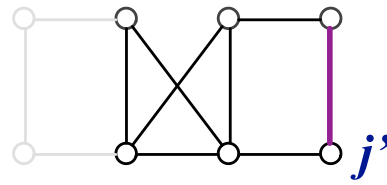
Figure 8: Network  $g$

The network  $g$  depicted in Figure 8 fails to satisfy  $k$ -WARPS with  $k = 2$ . In particular, we will show that  $(i, g) \sim_2 (i', g + i'j')$  and  $(j, g) \sim_2 (j', g + i'j')$ . The network in Figure 9 is the neighborhood of agent  $i$  in network  $g$ , while the network in Figure 10 shows the 2-neighborhood of agent  $i'$  in network  $g + i'j'$ . Figure 10 on the right is obtained by first adding a link between  $i'$  and  $j'$  (indicated by the color purple), and then removing the greyed out nodes and links.

Agents  $i$  and  $i'$  occupy similar positions in the respective neighborhoods, formally, there is an isomorphism  $\sigma$  of between the two neighborhood networks such that  $\sigma(i) = i'$ . Therefore we conclude that  $(i, g) \sim_2 (i', g + i'j')$ . Note that the network in Figure 10 is never observed in the data, but we postulate through  $\sim$  that agent  $i'$  preference in this network  $g + i'j'$ , is similar to agent  $i$ 's preference in the observed network  $g$ . Also observe that agents  $i, i'$  are never similar in the observed network  $g$ , as agent  $i$  has two links while  $i'$  has only one link in  $g$ .

Figure 9:  $(i, g)$ Figure 10:  $(i, g + i'j')$ 

The same reasoning shows that  $(j, g) \sim_2 (j', g + i'j')$ . Therefore the network  $g$  violates  $k$ -WARPS with  $k = 2$ , as we have found a pair of agents  $i, j$  and a non-link  $(i', j') \notin g$  such that  $(i, g) \sim_2 (i', g + i'j')$  and  $(j, g) \sim_2 (j', g + i'j')$ .

Figure 11:  $(j, g)$ Figure 12:  $(j', g + i'j')$ 

#### 7.4. Measure of Consistency with WARPS

To what extent is a network consistent with  $k$ -WARPS (Definition 7.3)? This section proposes a continuous measure of how consistent an observed network is to satisfying  $k$ -WARPS. When a network  $g$  violates WARPS, the interpretation is that there are links that should be dissolved or formed under pairwise stability. One way to quantify the inconsistency with WARPS is to look at the fraction of such agents who should dissolve some links under pairwise stability.

In Definition 7.5 below, we propose  $k$ -**stability** as a continuous measure of consistency with  $k$ -WARPS. The  $k$ -stability of a network is a real number between 0 and 1, and when  $k$ -stability of a network  $g$  is 1, the network  $g$  is fully consistent with  $k$ -WARPS. It has the intuitive appeal of decomposing consistency of a network with WARPS into agents whose behaviors are consistent with WARPS.

**Definition 7.5** (*k*-stability). The *k*-stability of a network  $g$ , is the fraction of agents in  $g$  that are *k*-stable. An agent  $i$  is *k*-stable if and only if there does not exist another agent  $j$  and a non-link  $i'j' \notin g$  such that  $(i, g) \sim_k (i', g + i'j')$  and  $(j, g) \sim_k (j', g + i'j')$

When an agent  $i$  is not *k*-stable, then inconsistency with pairwise stability can be resolved by having agent  $i$  dissolve a link. Therefore the measure of consistency with WARPS defined here intuitively measures the fraction of agents that should dissolve links.

#### 7.4.1. Precision and power of WARPS

One criticism of the measure in Definition 7.5 is that it does not take into account the precision or the power of WARPS. For some class of data, it may be that WARPS never rejects. For example, the empty and the complete networks always satisfy WARPS. Therefore, WARPS may have little precision or power for very sparse and dense networks. Indeed Section 9 shows that the average consistency with WARPS of a random (Erdős-Rényi) network varies considerably depending on the parameters  $n, p$ , where  $n$  is the number of nodes and  $p$  is the network density.

The second measure of consistency we will use is the **excess *k*-stability**. As described in Definition 7.6, it is the difference between the consistency of an observed network with WARPS, and the average consistency with WARPS of a suitably calibrated random network. This measure is the notion of predictive success proposed by Selten (1991), and used in the empirical revealed preference literature (Beatty and Crawford (2011)).<sup>9</sup>

**Definition 7.6** (Excess *k*-stability). The **excess *k*-stability** of a network  $g$  is calculated by subtracting the expected *k*-stability of an Erdős-Rényi random network<sup>10</sup> with the same parameters as  $g$ , from the *k*-stability of  $g$ .

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<sup>9</sup>Selten's index of predictive success is defined as the difference between the relative frequency of correct predictions (the 'hit rate') and the relative size of the set of predicted outcomes (the 'precision'). An attractive feature of Selten's predictive success is that it is uniquely characterized by a set of axioms.

<sup>10</sup>Specifically the  $G(n, p)$  random network where there are  $n$  nodes and each link is formed independently with probability  $p$

## 7.5. Discussion

This section discusses why this particular similarity relation is proposed (Definition 7.3). Firstly, it is a conservative notion of similarity that says agents have similar preferences only when their respective network structures are similar. The use of graph isomorphism allows preferences to depend freely on the network structures freely specifying how. More formally by definition of graph isomorphism, whenever  $(i, g) \sim_k (i', g')$ , we must have  $\theta_i(g) = \theta_{i'}(g')$ , as  $k \rightarrow \infty$ , and where  $\theta : N \times \mathcal{G} \rightarrow \mathbb{R}$  is any possible network characteristics.<sup>11</sup>

The definition of  $k$ -WARPS made it clear that as  $k \rightarrow \infty$ ,  $k$ -WARPS will never be rejected. For large  $k$ , pairwise stability has very little testable implication. In the empirical application that follows, we will set  $k = 2$ . That is, agents only care about network consisting of their friends, and their friends of friends. This is in line with the literature (de Paula et al. (2014); Mele (2013); Sheng (2014)).

In general, when we allow agents to have more heterogeneous preferences (through a less restrictive similarity relation  $\sim^{12}$ ), WARPS has less power. In Section 9, the power of WARPS is discussed, and it is defined as the probability that a random network would reject WARPS. We will see in Section 9, that the form of WARPS implemented in the empirical section (with  $k = 2$ ), has just enough power for our purpose.

## 8. EMPIRICAL APPLICATION

### 8.1. Data

In this section, I present an empirical application of  $k$ -stability, which measures consistency with  $k$ -WARPS (Definition 7.5). The data set in this section is taken from Banerjee et al. (2013). They collected detailed network data by surveying households about a wide range of interactions in 75 rural villages in India. This information was then used to create different network graphs

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<sup>11</sup>Formally, a network characteristics is defined to be a property that is preserved under types-preserving isomorphisms of networks. In other words, it is a structural property of the network itself, not of a specific drawing or labeling of the network. For instance,  $\theta_i(g)$  can be the network centrality of  $i$  in  $g$ , or the density of  $g$ .

<sup>12</sup>More formally,  $\sim$  is less restrictive than  $\sim'$  if  $\sim$  is a finer partition of  $N \times \mathcal{G}$  than  $\sim'$



for each village.<sup>13</sup> I will consider 4 different types of networked relationships. Note that agents in this context refers to households.

*Financial network* with respect to money: there is a financial relationship with regards to money between agents  $i$  and  $j$  if and only if agent  $i$  reports borrowing money from (or lending to) agent  $j$ , and agent  $j$  reports lending money to (or borrowing from)  $i$ . Similarly, I construct networks of financial relationships with regards to borrowing and lending rice or kerosene. *Advice network*: there is an advice relationship between agents  $i$  and  $j$  if and only if agent  $i$  reports receiving advice from (or giving to) agent  $j$ , and agent  $j$  reports giving advice to (or receiving from)  $i$ . *Social networks* are constructed in similar manner.<sup>14</sup>

Figure 13 shows the financial network of a particular village. This network does not satisfy  $k$ -WARPS with  $k = 2$ , and the degree to which it is consistent with 2-WARPS is measured at 80.8%, which is the 2-stability of the network (Definition 7.5). This is calculated as follows: 14 out of 73 households violate 2-WARPS agents (households), and these are the nodes whose links are bolded in the figure. Moreover, a (Erdős-Rényi) random network with the same parameter is expected to have a 2-stability score of 83.4%. The expected 2-stability of this network is then 80.8 minus 83.4.

## 8.2. Descriptive result

There is considerable variation in  $k$ -stability across villages for a given network type. For example, Figure 14 plots the histogram of  $k$ -stability with  $k = 2$  for the cross section of 75 village networks. Figure 15 shows how excess 2-stability varies across villages and network types. Figures 14 and 15 also highlight how the empirical distribution of  $k$ -stability can be very different from the empirical distribution of the corresponding excess  $k$ -stability.

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<sup>13</sup>More precisely, individuals are surveyed and a relationship between households exists if any household members indicated a relationship with members from the other household. Although surveyed individuals could name nonsurveyed individuals as friends, relatives, etc., such links are omitted unless both individuals were surveyed.

<sup>14</sup>*Social network*: there is a social relationship between agents  $i$  and  $j$  if and only if agent  $i$  reports visiting the house of (or receiving visits from) agent  $j$ , and agent  $j$  reports receiving visits from (or visiting the house of) agent  $i$ .

<sup>15</sup>There is a link between agents  $i$  and  $j$  if and only if agent  $i$  reports borrowing money from (or lending to) agent  $j$ , and agent  $j$  reports lending money to (or borrowing from)  $i$ .

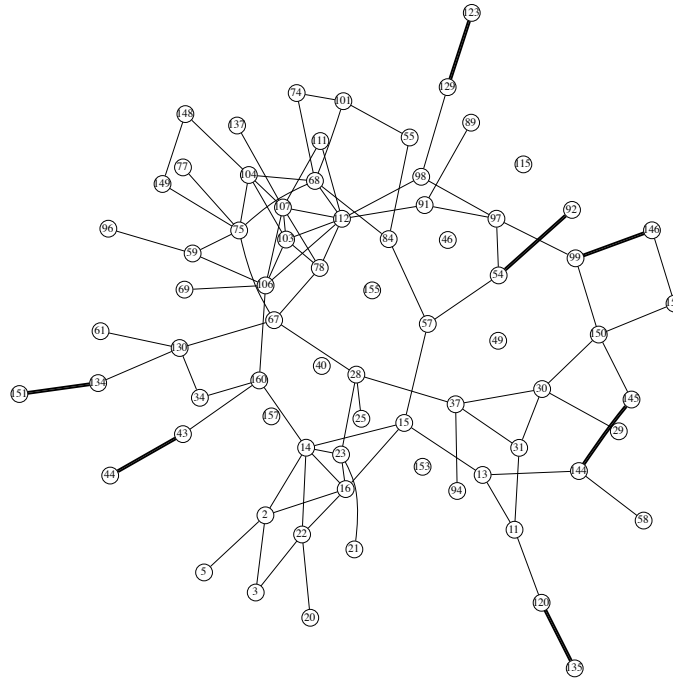


Figure 13: Network of borrowing and lending money<sup>15</sup> in village number 27 of Banerjee et al. (2013).

### 8.3. Correlates of $k$ -stability

This section examines the village level predictors of  $k$ -stability. Specifically, I will look at how the consistency of a village network with WARPS is related to the average characteristics of the village. In Table I, I provide results for an ordinary least squares (OLS) regression of consistency with WARPS as a function of various village characteristics. The measure of consistency with WARPS here is the excess 2-stability (Definition 7.6), averaged across financial, social and advice networks.

The main village characteristics that explain consistency with WARPS is the number of households and the fraction of individuals in the village who travel outside the village for work. In particular, they are both significantly

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Non-surveyed households are not included. Graph is produced using the spring algorithm.

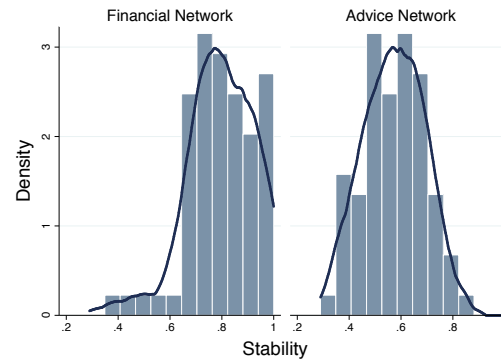


Figure 14: The histograms and kernel densities of 2-stability across 75 villages, for two main types of network: financial (money) and advice networks. The correlation between the two is 0.232. The histogram for 3-stability looks similar, but with mass shifted towards 1.

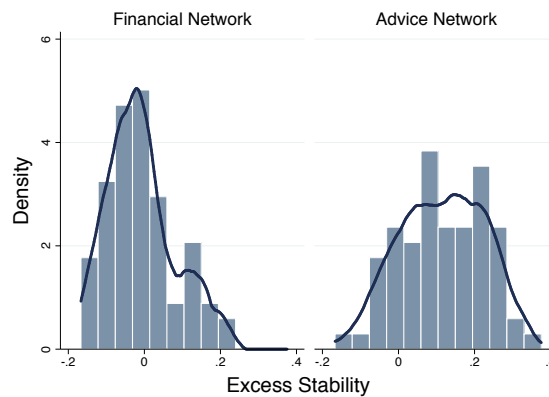


Figure 15: Comparing the excess 2-stability of various networks types across the 75 villages. The correlation between the two is  $-0.067$ .

negatively related to consistency with WARPS. While the first variable is readily available in the data set provided by [Banerjee et al. \(2013\)](#), the "Work Outside Village" variable is not. I construct this variable by calculating the fraction of respondents in a village who answer 'Yes' to the question 'Do you travel outside the village for work?'. This empirical relationship suggests that (i) larger village tends to be less consistent with WARPS, and (ii) village where individuals spend less time in the village tends to also be less

consistent with WARPS.

It is plausible that in a larger village, mutually beneficial relationships are not formed because agents have limited attention (Masatlioglu et al. (2012)) and do not consider all such beneficial relationships. This could be due to the geographical spread of the village as village size increases.

It is also possible that the variation in consistency with WARPS is driven by unequal measurement and sampling error across villages, where links are misreported with higher frequency in some villages.<sup>16</sup> Although households can name others not in the survey, I omit those links involving non-surveyed households (similar to Jackson et al. (2012)). Moreover, I construct the variable “Recip” that measures the proportion of reports that are reciprocated in a village, for instance, when household  $i$  reports borrowing from household  $j$ , how frequent does  $j$  also reports lending to  $i$ . This variable can be a proxy for measurement error in constructing village-level undirected network from survey data, but including it did not change the regression result in Table I. Finally, the consistency of financial network with WARPS has negligible correlation (even negative) with the stability of social network, even though the measures are calculated based on the same sampled nodes.

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<sup>16</sup>Chandrasekhar and Lewis (2011) examines the biases that would arise when sampled (and missing) network data are used in regression analysis.

Excess 2-Stability (Consistency with WARPS)					
Number of Households (per 1000)	-0.508** (0.216)	-0.530** (0.226)		-0.563** (0.219)	
Work Outside Village	-0.125* (0.0674)	-0.260** (0.120)		-0.292** (0.136)	-0.310** (0.134)
Average Education (per 100)		-1.227 (0.824)	-0.437 (0.703)	-1.453 (0.990)	-1.446 (1.001)
Average Age (per 100)		-0.912 (0.620)	-0.139 (0.523)	-1.168 (0.737)	-0.975 (0.728)
Fraction GM			0.0413 (0.0471)	-0.0103 (0.0400)	-0.0239 (0.0426)
Fraction Nonnatives			-0.0888 (0.0975)	-0.117 (0.119)	-0.102 (0.0987)
Average No. Rooms				0.0239 (0.0284)	0.0170 (0.0310)
Constant	0.0961*** (0.0226)	0.544* (0.282)	0.110 (0.208)	0.668** (0.326)	0.563* (0.319)
<i>N</i>	75	75	75	75	75
<i>R</i> <sup>2</sup>	0.092	0.144	0.021	0.163	0.101

Robust standard errors in parentheses

\*  $p < 0.10$ , \*\*  $p < 0.05$ , \*\*\*  $p < 0.01$

TABLE I

The relationship between consistency with WARPS and village characteristics. The consistency of village networks with WARPS is negatively correlated with the number of households and the fraction of villagers who travel outside the village for work. Excess stability takes into account the random variation in consistency with WARPS due to variation in network density and number of households. Other village-level variables I control for are the average education level; average age of the villagers; fraction belonging to the upper two castes; fraction who are not natives of the village; and the average number of rooms owned by each person (wealth indicator).

#### 8.4. Relation of $k$ -stability to diffusion of microfinance

In this section, I show how the consistency with WARPS is associated with economic outcomes, specifically, the diffusion of microfinance.<sup>17</sup>

<sup>17</sup>Potentially, the result obtained here is applicable for diffusion of information on networks in a broader sense. Galeotti et al. (2010); Jackson and Yariv (2007, 2010).

For each village, in addition to the network data, we also observe the nodes who were first informed of a new microfinance loan program, as well as the eventual adoption rate of microfinance. However, the full information is only available for 43 of the original 75 surveyed villages. The first-informed agents are also known as the injection point of the network, they are typically leaders in the village.

[Banerjee et al. \(2013\)](#) show that when the injection points occupy more central positions in the underlying social networks, then the eventual adoption rates of microfinance are greater. Using [Banerjee et al. \(2013\)](#)'s specification, I incorporate our measures of consistency with WARPS as additional variables. The main empirical finding is that when the underlying network is more consistent with WARPS, the positive relationship between network centrality of injection points and adoption rate is even more significant and stronger now.

Therefore, the consistency of WARPS is new network statistics that has potentially useful application: when a network is more consistent with WARPS, targeting of agents that have central positions in the network to increase the diffusion of information is more effective. Moreover the result is driven primarily by the consistency of advice network with WARPS. One expects the advice network to be the network type more salient for households passing on information about a new loan program.

The intuition for the result is as follows. Networks that are less consistent with WARPS are more difficult to rationalize as pairwise stable. There are only two reasons, either we do not expect pairwise stability to hold, or that there is unobserved heterogeneity not captured by the network structure. For either of these explanations, centrality of injection points *with respect to* the network structure does not correspond to the actual effectiveness of these agents in diffusion. Therefore, the relationship between network centrality and adoption rate is weaker as the underlying network violates WARPS more severely. This is precisely what we see in the data.

#### 8.4.1. Details of regression result

The first column of Table II replicates the regression result of [Banerjee et al. \(2013\)](#). The dependent variable is the microfinance take-up rate of non-leader households. The main explanatory variables introduced in [Banerjee et al. \(2013\)](#) are the communication and diffusion centrality of the injection

points, which are measures of how central individuals are in their social network with regard to spreading information. We can see in Table II that the diffusion centrality of injection points helps significantly in predicting eventual adoption of microfinance. The controls include Number of households, Savings, SHG participation, Fraction GM, and Fraction Leaders.<sup>18</sup>

In Column (2) and (3) of Table II, I introduced different measures of consistency with WARPS to the Banerjee et al. (2013)'s specification. Specifically, in Column (2) of Table II, a village's consistency with WARPS is defined to be the 2-stability (see Section 7.5), averaged across all the different network types: financial (money), financial (rice and kerosene), advice, and social networks. Analogously in Column (3), the measure of consistency is the excess 2-stability<sup>19</sup>, averaged across different network types.

As seen in Table II, our measures of consistency with WARPS added considerable explanatory power to the original model of diffusion. In all specifications, the interaction between consistency with WARPS and diffusion centrality is significantly positive. The main empirical finding is that when underlying networks are more consistent with WARPS, network centrality of injection points matters more for diffusion.

To get a sense of the magnitudes, using Column (3) of Table II, one standard deviation increase in diffusion centrality is associated with 6.7, 4.6, 2.1 percentage points increase in the dependent variable at 90th, 75th, 25th percentile of excess 2-stability. Magnitudes using Column (2) are similar. In comparison to Banerjee et al. (2013), one standard deviation increase in diffusion centrality is associated with an increase of 3.9 percentage points (the average participation rate is 18.5%)

Interestingly the result also suggests that network that violates WARPS more has higher diffusion. Consistency with WARPS by itself seems to have small but negative effect on adoption rate. One explanation is that random and non-network-based meetings of agents is important for the spread of information, and networks that are stable have less random meetings of agents.

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<sup>18</sup>Savings: Fraction of households engaging in formal savings; SHG participation: Fraction of households participating in a self-help group; Fraction GM: Fraction of households that are not from the scheduled castes and scheduled tribes; groups that historically have been relatively disadvantaged; Fraction Leaders: Fraction of households that are leaders

<sup>19</sup>The excess  $k$ -stability score is the  $k$ -stability score in excess of a randomly generated network with the same network parameters (see Definition 7.6).

Therefore our result suggests that for villages with higher stability, network-based diffusion is more prevalent; while for villages with lower stability, random non-network meetings drive diffusion.

#### 8.4.2. Robustness

In Table III, I repeat the same exercise but focusing instead on advice networks. Table III verifies that the result is not sensitive to different ways of defining consistency with WARPS. The interaction between consistency with WARPS and diffusion centrality is significantly positive in all 4 columns representing different ways of defining consistency with WARPS.

In Table IV, I present a placebo test of the main specification (Table II) using different network measures in place of consistency with WARPS. The network measures used are (1): Number of nodes (households); (2): Average degrees; (3) Average distance; (4) Average clustering; (5): Largest eigenvalue of the adjacency matrix. Column (6) shows that although average clustering has the same effect on the dependent variable as stability, this effect is not robust when communication centrality of injection points is used instead. Communication centrality is another measure of centrality that is similar to diffusion centrality (see Banerjee et al. (2013)).

Table V shows the correlation between our main measure of consistency and other commonly used network measures. Just like consistency with WARPS, these network measures are calculated using only information about the network structure, but nonetheless Table IV shows that, with the exception of the number of nodes (households), they were not significant in explaining diffusion.



	(1)	(2)	(3)
	Microfinance participation rate		
Diffusion Centrality	0.0222*** (0.00690)	-0.0886* (0.0470)	0.0165*** (0.00594)
Stability		-0.897*** (0.323)	
<b>Stability × Diffusion Centrality</b>		<b>0.143**</b> (0.0583)	
Excess Stability			-1.275*** (0.391)
<b>Excess Stability × Diffusion Centrality</b>			<b>0.230***</b> (0.0734)
Number of Households (per 1000)	-0.540*** (0.163)	-0.472*** (0.169)	-0.507*** (0.157)
Savings	-0.244* (0.138)	-0.205* (0.118)	-0.252* (0.138)
SHG Participation	-0.228 (0.200)	-0.179 (0.180)	-0.165 (0.199)
Fraction GM	-0.0354 (0.0396)	-0.0280 (0.0395)	-0.0410 (0.0404)
Fraction Leaders	0.270 (0.457)	0.301 (0.442)	0.222 (0.425)
<i>N</i>	43	43	43
<i>R</i> <sup>2</sup>	0.442	0.511	0.542

Robust standard errors in parentheses

\*  $p < 0.10$ , \*\*  $p < 0.05$ , \*\*\*  $p < 0.01$

TABLE II

The dependent variable is the microfinance take-up rate of non-leader households. The first column reproduces exactly the result of Banerjee et al. (2013). Two measures of consistency with WARPS are used: (excess) stability is the average (excess) 2-stability across all network types. Excess  $k$ -stability is the  $k$ -stability of a network in excess of a randomly generated network with the same parameter. Using Column (2), one standard deviation increase in diffusion centrality is associated with 6.8, 5.6, 3.0 percentage points increase in the dependent variable at 90th, 75th, 25th percentile of stability scores. Using Column (3), one standard deviation increase in diffusion centrality is associated with 6.7, 4.6, 2.1 percentage points increase in the dependent variable at 90th, 75th, 25th percentile of stability scores.

	(1)	(2)	(3)	(4)
Diffusion Centrality	-0.0527 (0.0367)	-0.133*** (0.0282)	0.00370 (0.00795)	0.0175*** (0.00548)
Stability	-0.760** (0.299)	-0.954*** (0.190)		
<b>Stability × Diffusion Centrality</b>	0.128* (0.0657)	0.200*** (0.0405)		
Excess Stability			-0.890*** (0.272)	-0.822*** (0.295)
<b>Excess Stability × Diffusion Centrality</b>			0.155** (0.0577)	0.188** (0.0707)
Constant	1.107*** (0.316)	1.191*** (0.281)	0.723** (0.280)	0.846*** (0.275)
<i>N</i>	43	43	43	43
<i>R</i> <sup>2</sup>	0.529	0.583	0.572	0.550

Robust standard errors in parentheses

\*  $p < 0.10$ , \*\*  $p < 0.05$ , \*\*\*  $p < 0.01$

TABLE III

The dependent variable is the microfinance take-up rate of non-leader households. Only Advice Networks are used. Demographic controls included as in Table II. In columns (1) and (2), stability is defined as the 2 and 3-stability of advice networks. In Columns (3) and (4), excess stability is defined as the excess 2 and 3-stability of advice networks. Magnitudes are comparable with Table II. For column (1), one standard deviation increase in the diffusion centrality is associated with 6.6, 5.3 and 1.8 percentage points increase in the dependent variable at 90th, 75th, 25th percentile of stability scores. For column (2), the numbers are 8.3, 5.8, and 1.2; for column (3): 7.2, 5.9, 1.3; for column (4): 8.1, 6.1, and 1.7

	(1)	(2)	(3)	(4)	(5)	(6)	(7)
	Microfinance participation rate						
Diffusion Centrality	0.0516*** (0.0142)	0.00486 (0.0515)	0.0318 (0.0724)	-0.0560*** (0.0162)	0.0506 (0.0429)		
Network Measures × Diffusion Centrality	-0.000349* (.000174)	0.00623 (0.0199)	-0.00219 (0.0156)	0.601*** (0.131)	-0.00699 (0.00971)		
Comm. Centrality						-0.695 (0.825)	0.449** (0.204)
Network Measures × Comm. Centrality						11.26 (6.725)	12.17*** (2.536)
Network Measures	.000556 (.000878)	-0.0739 (0.101)	0.0480 (0.0879)	-3.206*** (0.685)	-0.000738 (0.0524)	-0.671 (0.406)	-0.847*** (0.266)
Controls	Yes	Yes	Yes	Yes	Yes	Yes	Yes
<i>N</i>	43	43	43	43	43	43	43
<i>R</i> <sup>2</sup>	0.492	0.470	0.468	0.582	0.495	0.443	0.574

Robust standard errors in parentheses

\*  $p < 0.10$ , \*\*  $p < 0.05$ , \*\*\*  $p < 0.01$

TABLE IV

Placebo test of the main specification (Table II) using different network measures in place of consistency with WARPS. The network measures used are (1): Number of nodes (households); (2): Average degrees; (3) Average distance; (4) Average clustering; (5): Largest eigenvalue of the adjacency matrix. Column (6) shows that although average clustering has the same effect on the dependent variable as stability, this effect is not robust when communication centrality of injection points is used instead. Communication centrality is another measure of centrality that is similar to diffusion centrality (see Banerjee et al. (2013)). Finally, Column (7) shows that our new network measure, excess 2-stability, is robust to using communication centrality in place of diffusion centrality.

	(1)	(2)	(3)	(4)	(5)	(6)	(7)
	Excess Stability						
No. of Households (per 1000)	-0.549** (0.217)						0.0564 (0.464)
Average degrees		0.0161 (0.0217)					-0.0494 (0.0372)
Average distance			-0.0252* (0.0146)				-0.0187 (0.0207)
Average clustering				0.415*** (0.151)			0.568*** (0.199)
1st Eigenvalue					0.00718 (0.0124)		0.00926 (0.0249)
2nd Eig. Stoch. Adj						-0.620 (0.610)	-0.958 (0.644)
Constant	0.0732*** (0.0202)	-0.0203 (0.0581)	0.140** (0.0659)	-0.0426* (0.0246)	-0.00969 (0.0560)	0.634 (0.601)	1.050 (0.697)
<i>N</i>	75	75	75	75	75	75	75
<i>R</i> <sup>2</sup>	0.062	0.012	0.052	0.121	0.005	0.016	0.192

Standard errors in parentheses

\*  $p < 0.10$ , \*\*  $p < 0.05$ , \*\*\*  $p < 0.01$

TABLE V

Correlation of excess 2-stability with other network measures, at the village-level. All measures are averaged over the four types of networks. These network measures are the ones used in [Banerjee et al. \(2013\)](#)

## 9. CONSISTENCY OF A RANDOM NETWORK WITH WARPS

In this section, I investigate the consistency with WARPS of a random network. The likelihood that a random network violates WARPS will be our measure of **power**. Moreover, the result obtained here will be used as a benchmark for gauging how severe real-world networks violate WARPS (see Excess  $k$ -stability of Definition 7.6).

Using  $k$ -stability (Definition 7.5) as our measure of consistency with WARPS, I will examine the  $k$ -stability of the canonical Erdős-Rényi random graph model. Here, I consider the  $G(n, p, \lambda)$  random graph model, where a network with  $n$  number of nodes is randomly generated as follows: each link is included in the network with probability  $p$  (independently). With probability  $\lambda$ , a node is assigned as type 0, and with probability  $1 - \lambda$ , a node is assigned as type 1.

We are interested in the **expected stability function**  $E(n, p, \lambda; k)$ , which tells us on average, what is the  $k$ -stability for a network randomly drawn from the random graph model,  $G(n, p, \lambda)$ . For instance, if  $E(n_0, p_0, \lambda_0) = 0.4$ , then we expect a network randomly generated from  $G(n_0, p_0, \lambda_0)$  to have a  $k$ -stability score of 0.4. Therefore the score of 0.4 is the appropriate benchmark for an observed network with  $n_0$  number of nodes, density of  $p_0$ , and  $\lambda_0$  proportion of type-0. In this section, we will focus on  $k$ -stability with  $k = 2, 3$ , but the same analysis can be carried out for different  $k$ -stability scores.

**Definition 9.1.** Let  $\mathcal{S}(g; k)$  be the  $k$ -stability of the network  $g$  as defined in 7.4. Let  $n, p, \lambda$  be the number of nodes, the density, and the proportion of type-0 in the network  $g$  respectively. The **excess  $k$ -stability** of a network  $g$  is  $\mathcal{S}(g; k) - E(n, p, \lambda; k)$

First I fix  $\lambda = 0$  (or equivalently  $\lambda = 1$ ), that is, assuming that nodes do not have different types or characteristics. Although I am not able to analytically characterize the expected stability function, Figure 16 shows simulation evidence that the expected  $k$ -stability of a random network depends only on the product  $(n - 1)p$ , which is the average degree of the network, i.e. how many links a node has on average. It is known that  $(n - 1)p$  characterized certain phase transitions of Erdős-Rényi random networks.<sup>20</sup>

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<sup>20</sup>For instance, when the average degree is less than one, the network consists of many components that are small relative to the number of nodes, and when the average degree

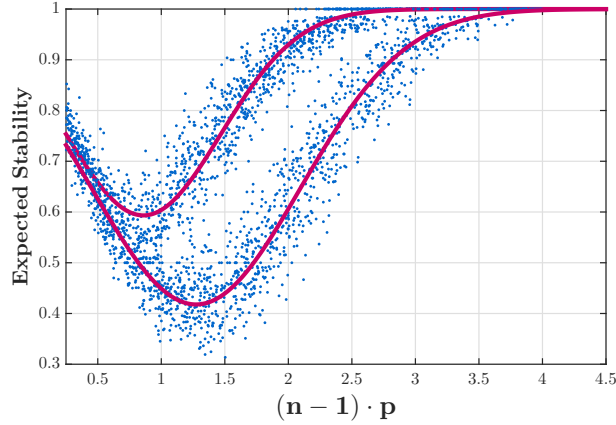


Figure 16: The simulated expected 2 and 3-stability can be accurately described by an inverted Gaussian function that depends only on the average degree  $(n-1)p$ . That is, the curves above are obtained by fitting  $1 - \alpha \cdot \exp \left[ -\left( \frac{(n-1)p - \beta}{\gamma} \right)^2 \right]$  to the simulation points. The curve for expected 3-stability lies strictly on top.

Moreover as can be seen in Figure 16, the simulated expected stability can be described by an inverted Gaussian function remarkably well. That is, 
$$E(n, p) = 1 - \alpha \cdot \exp \left[ -\left( \frac{(n-1)p - \beta}{\gamma} \right)^2 \right].$$

Ultimately we wish to know the general shape of the expected stability function, with  $\lambda \in (0, 0.5]$ . In Figure 17, I plot a non-parametric smoothed surface over the simulated expected stability. This is obtained by first simulating the expected stability over a range of values for  $n$ ,  $p$ , and  $\lambda$ ; and then secondly, computing the local, non-parametric regression known as the LOESS Curve.

Again the non-parametric plot of the simulated expected stability in Figure 17 is suggestive of an inverted multivariate Gaussian function. That is, I obtain an  $R^2$  of 0.96 when I fit to the simulated data, the inverted Gaussian function of the form  $E(n, p, \lambda) = 1 - \alpha \exp \left[ -\left( \frac{(n-1)p - \beta}{\gamma} \right)^2 \right]$ , where the height and location of the peak ( $\alpha$  and  $\beta$  respectively) is a linear function of

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is greater than one, there is a unique giant component.

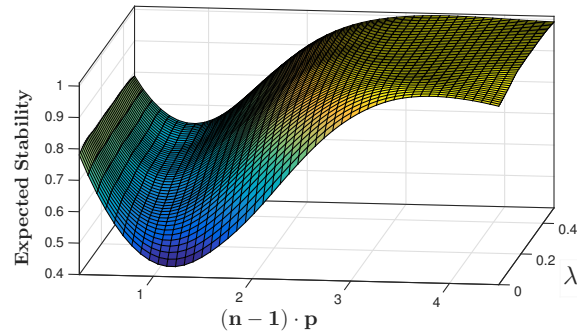


Figure 17: Non-parametric smoothed surface (Loess Curve) of the simulated expected stability function  $E(n, p, \lambda)$ . The function  $E(n, p, \lambda)$  is the expected 2-stability score of a network randomly drawn from  $G(n, p, \lambda)$ . Stability is expected to be lowest around  $(n - 1)p = 1$ , which corresponds to the well-known phase transition for the rise of a giant component in Erdős-Rényi random networks. As  $\lambda$  increases, stability is expected to increase.

$\lambda$ .<sup>21</sup>

Figure 16 is obtained by first drawing 1000 pairs of  $(n, p)$  uniformly from  $n \in \{40, \dots, 70\}$ , and  $p$  such that  $np \in [0, 4.5]$ , then for each parameter  $(n, p)$ , 10 random networks are drawn from  $G(n, p)$  for a total of 10,000 networks. The 2 and 3-stability scores are then computed for each of these networks. To obtain Figure 17, I simulate the expected stability function at 10,000 distinct values of  $(n, p, \lambda)$  drawn uniformly from  $n \in \{40, \dots, 70\}$ ,  $\lambda \in [0, 0.5]$ , and  $p$  such that  $np \in [0, 4.5]$ . For each  $(n, p, \lambda)$ , I draw 10 random networks from  $G(n, p, \lambda)$ , and I then compute the 2-stability score for each randomly drawn networks (100,000 in total).

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<sup>21</sup>There is however a caveat: the Gaussian function appears only to match the non-parametric surface when  $\lambda$  is not too close to the boundary of 0 and 0.5, and  $(n - 1)p$  not too close to the boundary of 0.

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